# Probabilistic Graphical Models 

Exact Inference

Complexity of VE
Belief Propagation

## Summary

- Simple algorithm
- Works for both BNs and MNs
- Factor product and summation steps can be done in any order, subject to:
- when $Z$ is eliminated, all factors involving $Z$ have been multiplied in


## Variable Elimination

- Goal: $\quad P(J)$
- Eliminate: $C, D, I, H, G, L, S$
$\sum_{C, D, I, G, S, L, H} \phi_{C}(C) \phi_{D}(C, D) \phi_{I}(I) \phi_{G}(G, I, D) \phi_{S}(S, I) \phi_{L}(L, G) \phi_{J}(J, L, S) \phi_{H}(H, G, J)$



## Complexity

$$
\begin{aligned}
\psi_{k}\left(\boldsymbol{X}_{k}\right) & =\prod_{i=1}^{m_{k}} \phi_{i} \\
\tau_{k}\left(\boldsymbol{X}_{k} \backslash\{Z\}\right) & =\sum_{Z} \psi_{k}\left(\boldsymbol{X}_{k}\right)
\end{aligned}
$$

## Factor product of $m_{k}$ factors

Marginalization of a variable

## Factor Product

$$
\begin{gathered}
N_{k}=\left|\operatorname{Val}\left(\boldsymbol{X}_{\boldsymbol{k}}\right)\right| \\
N_{k} \text { rows }
\end{gathered}
$$

| $a^{1}$ | $b^{1}$ | 0.5 |
| :---: | :---: | :---: |
| $a^{1}$ | $b^{2}$ | 0.8 |
| $a^{2}$ | $b^{1}$ | 0.1 |
| $a^{2}$ | $b^{2}$ | 0 |
| $a^{3}$ | $b^{1}$ | 0.3 |
| $a^{3}$ | $b^{2}$ | 0.9 |


| $b^{1}$ | $c^{1}$ | 0.5 |
| :--- | :--- | :--- |
| $b^{1}$ | $c^{2}$ | 0.7 |
| $b^{2}$ | $c^{1}$ | 0.1 |
| $b^{2}$ | $c^{2}$ | 0.2 |

$$
\psi_{k}\left(\boldsymbol{X}_{k}\right)=\prod_{i=1}^{m_{k}} \phi_{i}
$$

Cost: $\left(m_{k}-1\right) \times N_{k}$

| $a^{1}$ | $b^{1}$ | $c^{1}$ | $0.5 \cdot 0.5=0.25$ |
| :--- | :--- | :--- | :---: |
| $a^{1}$ | $b^{1}$ | $c^{2}$ | $0.5 \cdot 0.7=0.35$ |
| $a^{1}$ | $b^{2}$ | $c^{1}$ | $0.8 \cdot 0.1=0.08$ |
| $a^{1}$ | $b^{2}$ | $c^{2}$ | $0.8 \cdot 0.2=0.16$ |
| $a^{2}$ | $b^{1}$ | $c^{1}$ | $0.1 \cdot 0.5=0.05$ |
| $a^{2}$ | $b^{1}$ | $c^{2}$ | $0.1 \cdot 0.7=0.07$ |
| $a^{2}$ | $b^{2}$ | $c^{1}$ | $0 \cdot 0.1=0$ |
| $a^{2}$ | $b^{2}$ | $c^{2}$ | $0 \cdot 0.2=0$ |
| $a^{3}$ | $b^{1}$ | $c^{1}$ | $0.3 \cdot 0.5=0.15$ |
| $a^{3}$ | $b^{1}$ | $c^{2}$ | $0.3 \cdot 0.7=0.21$ |
| $a^{3}$ | $b^{2}$ | $c^{1}$ | $0.9 \cdot 0.1=0.09$ |
| $a^{3}$ | $b^{2}$ | $c^{2}$ | $0.9 \cdot 0.2=0.18$ |

## Factor Marginalization

| $a^{1}$ | $b^{1}$ | $c^{1}$ | $0.5 \cdot 0.5=0.25$ |
| :---: | :---: | :---: | :---: |
| $a^{1}$ | $b^{1}$ | $c^{2}$ | $0.5 \cdot 0.7=0.35$ |
| $a^{1}$ | $b^{2}$ | $c^{1}$ | $0.8 \cdot 0.1=0.08$ |
| $a^{1}$ | $b^{2}$ | $c^{2}$ | $0.8 \cdot 0.2=0.16$ |
| $a^{2}$ | $b^{1}$ | $c^{1}$ | $0.1 \cdot 0.5=0.05$ |
| $a^{2}$ | $b^{1}$ | $c^{2}$ | $0.1 \cdot 0.7=0.07$ |
| $a^{2}$ | $b^{2}$ | $c^{1}$ | $0 \cdot 0.1=0$ |
| $a^{2}$ | $b^{2}$ | $c^{2}$ | $0 \cdot 0.2=0$ |
| $a^{3}$ | $b^{1}$ | $c^{1}$ | $0.3 \cdot 0.5=0.15$ |
| $a^{3}$ | $b^{1}$ | $c^{2}$ | $0.3 \cdot 0.7=0.21$ |
| $a^{3}$ | $b^{2}$ | $c^{1}$ | $0.9 \cdot 0.1=0.09$ |
| $a^{3}$ | $b^{2}$ | $c^{2}$ | $0.9 \cdot 0.2=0.18$ |

$$
\begin{aligned}
& \phi\left(a^{1}, c^{1}\right)=\sum_{b} \phi\left(a^{1}, c^{1}, b\right) \\
& \qquad \begin{array}{|c|c|c|}
\hline a^{1} & c^{1} & 033 \\
\hline a^{1} & c^{2} & 0.51 \\
\hline a^{2} & c^{1} & 0.05 \\
\hline a^{2} & c^{2} & 0.07 \\
\hline a^{3} & c^{1} & 0.24 \\
\hline a^{3} & c^{2} & 0.39 \\
\hline
\end{array}
\end{aligned}
$$

Cost: $\sim N_{k}$ additions

## Complexity

- Complexity of Variable Elimination
- Start with $m$ factors
- $m \leq n$ for Bayesian networks (one for every variable)
- can be larger for Markov networks
- At each elimination step generate 1 factor,
- At most n elimination steps
- Total number of factors: $m^{\star} \leqslant m+n$


## Complexity

- $N=\max \left(N_{k}\right)=$ size of the largest factor
- Number of product operations:
- $\quad \sum_{k}\left(m_{k}-1\right) * N_{k} \leq \sum_{k}\left(m_{k}-1\right) * N=$

$$
=N * \sum_{k}\left(m_{k}-1\right) \leq N *(m+n)
$$

- Number of sum operations: $n * \sum_{k} N_{k} \leq n * N$
- Linear in $N$ and $m^{*}$


## Complexity

- $N=\max \left(N_{k}\right)=$ size of the largest factor
- Linear in $N$ and $m^{*}$
- But!
- $N_{k}=\left|\operatorname{Val}\left(\boldsymbol{X}_{k}\right)\right|=O\left(d^{r_{k}}\right)$
- $d=\max \left(\left|\operatorname{Val}\left(\boldsymbol{X}_{\boldsymbol{i}}\right)\right|\right) \mathrm{d}$ values in their scope
- $r_{k}=\left|\mathbf{X}_{k}\right|=$ cardinality of the scope of the $k$-th factor


## Elimination ordering?

| Step | Variable <br> eliminated | Factors <br> used | Variables <br> involved | New <br> factor |
| :---: | :---: | :---: | :---: | :---: |
| 1 | $C$ | $\phi_{C}(C), \phi_{D}(D, C)$ | $D, C$ | $\tau_{1}(D)$ |
| 2 | $D$ | $\tau_{1}(D), \phi_{G}(G, I, D)$ | $G, I, D$ | $\tau_{2}(G, I)$ |
| 3 | $H$ | $\tau_{2}(G, I), \phi_{I}(I), \phi_{S}(S, I)$ | $H, I, S$ | $\tau_{3}(G, S)$ |
| 4 | $G$ | $\phi_{H}(H, G, J)$ | $G, J, L, S$ | $\tau_{4}(G, J)$ |
| 5 | $S$ | $\tau_{3}(G, S), \tau_{4}(G, J) \phi_{L}(L, G)$ | $\tau_{5}(J, L, S)$ |  |
| 6 | $L$ | $\tau_{5}(J, L, S) \phi_{J}(J, L, S)$ | $J, L$ | $\tau_{6}(J, L)$ |
| 7 | $\tau_{6}(J, L)$ | $\tau_{7}(J)$ |  |  |

## Different elimination ordering?

| Step | Variable <br> eliminated |
| :---: | :---: |
| 1 | $G$ |
| 2 | $I$ |
| 3 | $S$ |
| 4 | $H$ |
| 5 | $C$ |
| 7 | $D$ |


$\sum_{C, D, I, G, S, L, H} \phi_{C}(C) \phi_{D}(C, D) \phi_{I}(I) \phi_{G}(G, I, D) \phi_{S}(S, I) \phi_{L}(L, G) \phi_{J}(J, L, S) \phi_{H}(H, G, J)$

## Different elimination ordering?

| Step | Variable <br> eliminated | Factors <br> used | Variables <br> involved | New <br> factor |
| :---: | :---: | :---: | :---: | :---: |
| 1 | $G$ | $\phi_{G}(G, I, D), \phi_{L}(L, G), \phi_{H}(H, G, J)$ | $G, I, D, L, J, H$ | $\tau_{1}(I, D, L, J, H)$ |
| 2 | $I$ | $\phi_{I}(I), \phi_{S}(S, I), \tau_{1}(I, D, L, S, J, H)$ | $S, I, D, L, J, H$ | $\tau_{2}(D, L, S, J, H)$ |
| 3 | $S$ | $\phi_{J}(J, L, S), \tau_{2}(D, L, S, J, H)$ | $D, L, S, J, H$ | $\tau_{3}(D, L, J, H)$ |
| 4 | $L$ | $\tau_{3}(D, L, J, H)$ | $D, L, J, H$ | $\tau_{4}(D, J, H)$ |
| 5 | $H$ | $\tau_{4}(D, J, H)$ | $D, J, H$ | $\tau_{5}(D, J)$ |
| 6 | $C$ | $\phi_{C}(C), \phi_{D}(D, C)$ | $D, J, C$ | $\tau_{6}(D)$ |
| 7 | $D$ | $\tau_{5}(D, J), \tau_{6}(D)$ | $D, J$ | $\tau_{7}(J)$ |

## Complexity and Elimination Ordering

Eliminate A first:

Eliminate $B_{i}$ 's first:


## Graphical Perspective

$$
\sum_{C, D, I, G, S, L, H} \phi_{C}(C) \phi_{D}(C, D) \phi_{I}(I) \phi_{G}(G, I, D) \phi_{S}(S, I) \phi_{L}(L, G) \phi_{J}(J, L, S) \phi_{H}(H, G, J)
$$



## Step 0: Moralize the Graph

$$
\sum_{C, D, I, G, S, L, H} \phi_{C}(C) \phi_{D}(C, D) \phi_{I}(I) \phi_{G}(G, I, D) \phi_{S}(S, I) \phi_{L}(L, G) \phi_{J}(J, L, S) \phi_{H}(H, G, J)
$$



## Step 1: Eliminate C

| Step | Variable <br> eliminated | Variables <br> involved | New <br> factor |
| :---: | :---: | :---: | :---: |
| 1 | $C$ | $D, C$ | $\tau_{1}(D)$ |
| 2 | $D$ | $G, I, D$ | $\tau_{2}(G, I)$ |
| 3 | $I$ | $G, I, S$ | $\tau_{3}(G, S)$ |
| 4 | $G$ | $H, G, J$ | $\tau_{4}(G, J)$ |
| 5 | $S$ | $J, L, S$ | $\tau_{5}(J, L, S)$ |
| 6 | $L$ | $J, L$ | $\tau_{6}(J, L)$ |
| 7 |  | $\tau_{7}(J)$ |  |



Step 2: Eliminate D

| Step | Variable <br> eliminated | Variables <br> involved | New <br> factor |
| :---: | :---: | :---: | :---: |
| 1 | $C$ | $D, C$ | $\tau_{1}(D)$ |
| 2 | $D$ | $G, I, D$ | $\tau_{2}(G, I)$ |
| 3 | $I$ | $G, I, S$ | $\tau_{3}(G, S)$ |
| 4 | $G$ | $H, G, J$ | $\tau_{4}(G, J)$ |
| 5 | $S$ | $J, L, S$ | $\tau_{5}(J, L, S)$ |
| 6 | $L$ | $J, L$ | $\tau_{6}(J, L)$ |
| 7 |  |  | $\tau_{7}(J)$ |



## Step 3: Eliminate I

| Step | Variable <br> eliminated | Variables <br> involved | New <br> factor |
| :---: | :---: | :---: | :---: |
| 1 | $C$ | $D, C$ | $\tau_{1}(D)$ |
| 2 | $D$ | $G, I, D$ | $\tau_{2}(G, I)$ |
| 3 | $I$ | $G, I, S$ | $\tau_{3}(G, S)$ |
| 4 | $G$ | $H, G, J$ | $\tau_{4}(G, J)$ |
| 5 | $S$ | $J, L, S$ | $\tau_{5}(J, L, S)$ |
| 6 | $L$ | $J, L$ | $\tau_{6}(J, L)$ |
| 7 |  |  | $\tau_{7}(J)$ |



## Step 4: Eliminate H

| Step | Variable <br> eliminated | Variables <br> involved | New <br> factor |
| :---: | :---: | :---: | :---: |
| 1 | $C$ | $D, C$ | $\tau_{1}(D)$ |
| 2 | $D$ | $G, I, D$ | $\tau_{2}(G, I)$ |
| 3 | $I$ | $G, I, S$ | $\tau_{3}(G, S)$ |
| 4 | $G$ | $H, G, J$ | $\tau_{4}(G, J)$ |
| 5 | $S$ | $J, L, S$ | $\tau_{5}(J, L, S)$ |
| 6 | $L$ | $J, L$ | $\tau_{6}(J, L)$ |
| 7 |  |  | $\tau_{7}(J)$ |



Steps 5,6,7: Eliminate G

| Step | Variable <br> eliminated | Variables <br> involved | New <br> factor |
| :---: | :---: | :---: | :---: |
| 1 | $C$ | $D, C$ | $\tau_{1}(D)$ |
| 2 | $D$ | $G, I, D$ | $\tau_{2}(G, I)$ |
| 3 | $I$ | $G, I, S$ | $\tau_{3}(G, S)$ |
| 4 | $G$ | $H, G, J$ | $\tau_{4}(G, J)$ |
| 5 | $S$ | $J, L, S$ | $\tau_{5}(J, L, S)$ |
| 6 | $L$ | $J, L$ | $\tau_{6}(J, L)$ |
| 7 |  |  | $\tau_{7}(J)$ |



S

## Induced Graph of an elimination ordering

- The induced graph $I_{\Phi, a}$ over factors $\Phi$ and ordering $a$ :
- Undirected graph
- $X_{i}$ and $X_{j}$ are connected if they appeared in the same factor in a run of the VE algorithm using $a$ as the ordering

Theorem: Every factor produced during VE is a clique in the induced graph

$$
\begin{gathered}
\tau_{1}(D)=\sum_{C} \phi_{C}(C) \phi_{D}(C, D) \\
\tau_{2}(G, I)=\sum_{D} \phi_{G}(G, I, D) \tau_{1}(D) \\
\tau_{3}(S, G)=\sum_{I} \phi_{S}(S, I) \phi_{I}(I) \tau_{2}(G, I) \\
\tau_{4}(G, J)=\sum_{H} \phi_{H}(H, G, J) \\
\tau_{5}(L, J)=\sum_{G} \frac{\phi_{L}(L, G) \tau_{3}(S, G) \tau_{4}(G, J)}{\tau_{6}=\sum_{L, S} \phi_{J}(J, L, S) \tau_{5}(L, J)}
\end{gathered}
$$

## Induced Graph of an elimination ordering

- The induced graph $I_{\Phi, a}$ over factors $\Phi$ and ordering $a$ :
- Undirected graph
- $X_{i}$ and $X_{j}$ are connected if they appeared in the same factor in a run of the VE algorithm using $a$ as the ordering

Theorem: Every maximal clique in the graph is a factor produced during VE Consider a maximal clique some variable is first to be eliminated once a variable is eliminated:

no new neighbor $\Rightarrow$ when eliminated it already had all the clique members as neighbors
$\Rightarrow$ participated in factors with all the other variables
$\Rightarrow$ when multiplied together, we have a factor oven all of them

## Complexity based on graphs

- The width of an induced graph is the number of nodes in the largest clique in the graph minus 1
- Minimal induced width of a graph K is $\min _{a}\left(\operatorname{width}\left(I_{K, a}\right)\right)$

- Provides a lower bound on best performance of VE to a model factorizing over K


## Finding a good elimination ordering

Theorem: For a graph H, determining whether there exists an elimination ordering for H with induced width K is NPcomplete

Note: This NP-hardness result is distinct from the NPhardness result of inference

- Even given the optimal ordering, inference may still be exponential


## Finding a good elimination ordering

- Greedy search using heuristic cost function - At each point, eliminate node with smallest cost
- Possible cost functions:
- min-neighbors: \# neighbors in current graph
- min-weight: weight (\# values) of factor formed
- min-fill: number of new fill edges
- weighted min-fill: total weight of new fill edges (edge weight = product of weights of the 2 nodes)


## Finding a good elimination ordering

Finding Elimination Orderings

- Theorem: The induced graph is triangulated
- No loops of length $>3$ without a "bridge"
-all cycles of four or more vertices have a chord, which is an edge that is not part of the cycle but connects two vertices of the cycle.
- Can find elimination ordering by finding a low-width triangulation of original graph $\mathrm{H}_{\Phi}$


## Example: Robot localization

## Robot Localization \& Mapping



## Example: Robot localization

Square Root SAM, F. Dellaert and M. Kaess, IJRR, 2006
Robot Localization \& Mapping


## Example: Robot localization

Square Root SAM, F. Dellaert and M. Kaess, IJRR, 2006
Eliminate Poses then Landmarks
Induced graph

## Example: Robot localization

Square Root SAM, F. Dellaert and M. Kaess, IJRR, 2006
Eliminate Landmarks then Poses


## Example: Robot localization

## Min-Fill Elimination



## Summary

Variable elimination allows computation of marginals / conditionals

Algorithm is valid for any graphical model
Suffices to show variable elimination for MRFs, since Bayes nets $\rightarrow$ MRFs by moralization

Worst-case complexity is dependent on elimination order, and is exponential in number of variables

Optimal ordering = treewidth, is NP-hard to compute

## Pt. 2: Message-passing inference

Variable Elimination

| Step | Variable <br> eliminated | Factors <br> used | Variables <br> involved | New <br> factor |
| :---: | :---: | :---: | :---: | :---: |
| 1 | $C$ | $\phi_{C}(C), \phi_{D}(D, C)$ | $D, C$ | $\tau_{1}(D)$ |
| 2 | $I$ | $\tau_{1}(D), \phi_{G}(G, I, D)$ | $D, I, G$ | $\tau_{2}(G, I)$ |
| 3 | $H$ | $\tau_{2}(G, I), \phi_{I}(I), \phi_{S}(S, I)$ | $G, I, S$ | $\tau_{3}(G, S)$ |
| 4 | $G$ | $\phi_{H}(H, G, J)$ | $G, J, L, S$ | $\tau_{4}(G, J)$ |
| 5 | $S$ | $\tau_{3}(G, S), \tau_{4}(G, J) \phi_{L}(L, G)$ | $\tau_{5}(J, L, S)$ |  |
| 6 | $L$ | $\tau_{5}(J, L, S) \phi_{J}(J, L, S)$ | $J, L$ | $\tau_{6}(J, L)$ |
| 7 | $\tau_{6}(J, S)$ | $\tau_{7}(J)$ |  |  |

## Understanding Variable Elimination



## Cluster Trees

Clusters


Variable Elimination as message passing

Message passing


$$
\begin{array}{ll}
\delta_{12}(D)=\sum_{D} \phi(C) \phi(C, D) & \delta_{35}(G, S)=\sum_{I} \phi(G, I, S) \delta_{35}(G, I) \\
\delta_{23}(G, I)=\sum_{D} \phi(D, I, G) \delta_{12}(D)
\end{array}
$$

Message passing

$$
\begin{aligned}
\delta_{12}(D) & =\sum_{D} \phi(C) \phi(C, D) \\
\delta_{23}(G, I) & =\sum_{D} \phi(D, I, G) \delta_{12}(D) \\
\delta_{35}(G, S) & =\sum_{I} \phi(G, I, S) \delta_{35}(G, I) \\
\delta_{45}(G, J) & =\sum_{H} \phi(G, H, J) \\
\delta_{56}(J, S, L) & =\sum_{G} \delta_{45}(G, J) \delta_{35}(G, S) \phi(L, G)
\end{aligned}
$$

## Clique-Tree Message Passing

1. Pick a node to be your root.
2. For each node $i$, initialize the potential of the node

$$
\psi_{i}=\prod_{i} \phi_{i}
$$

3. Start from a leaf and send message to all neighbors

$$
\delta_{i \rightarrow j}=\sum c_{i}-s_{i, j} \psi_{i} \cdot \prod_{k \in\left(\mathrm{Nb}_{i}-\{j\}\right)} \delta_{k \rightarrow i}
$$

4. Repeat for every node that is ready to transmit a message (i.e., has received messages from every neighbor)

## Properties of Cluster Trees

## Family Preservation:

For each factor $\phi_{k} \in \Phi$, there exists a cluster $C_{i}$ s.t. Scope $\left[\phi_{k}\right] \subseteq \boldsymbol{C}_{i}$ every factor has a node that can accommodate it

Running Intersection:
For each pair of clusters $C_{\mathrm{i}}, C_{\mathrm{j}}$ and variable $X \in C_{i} \cap C_{j}$, in the unique path between $C_{i}$ and $C_{j}$, all clusters and sepsets contain $X$. clusters that include the same variable need to communicate for consistency

## Running Intersection



Which clusters need to include X?

